

Transparent Bridging and VLAN

Plug and Play Networking

Algorhyme



*I think that I shall never see
a graph more lovely than a tree
a graph whose crucial property
is loop-free connectivity.*

*A tree which must be sure to span
so packets can reach every lan.
first the root must be selected
by ID it is elected.*

*least cost paths to root are traced,
and in the tree these paths are place.
mesh is made by folks like me;
bridges find a spanning tree.*



- Bridges came **after** routers!
- First bridge designed by **Radia Perlman**
 - ◆ Ethernet has size limitations
 - ◆ Routers were single protocol and expensive
- Spanning Tree because Ethernet had no hop count
- IEEE 802.1D

What is Bridging?

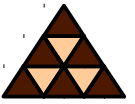


- **Layer 2 packet forwarding principle**
- **Separate two (or more) shared-media LAN segments with a bridge**
 - ◆ Only frames destined to the other LAN segment are forwarded
 - ◆ **Number of collisions reduced (!)**
- **Different bridging principles**
 - ◆ Ethernet: **Transparent Bridging**
 - ◆ Token Ring: **Source Route Bridging**

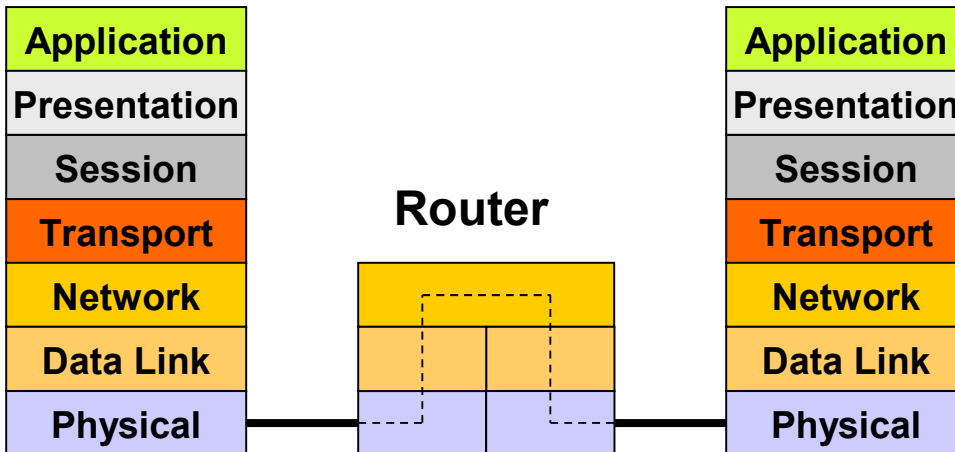
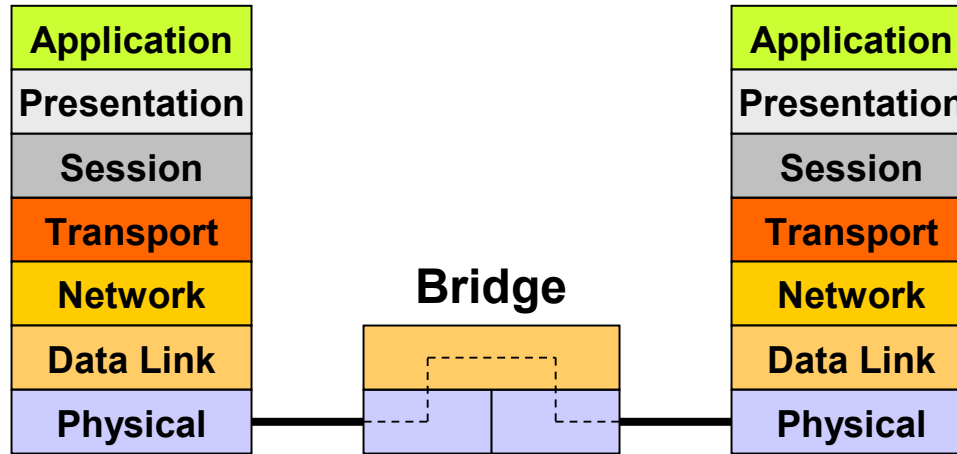
Bridging vs Routing



- **Bridging works on OSI layer 2**
 - ◆ Forwarding of **frames**
 - ◆ Use **MAC** addresses only
 - ◆ Termination of physical layer (!)
- **Routing works on OSI layer 3**
 - ◆ Forwarding of **packets**
 - ◆ Use **routable** addresses only (e.g. IP)
 - ◆ Termination of both layer 1 and 2



OSI Comparison

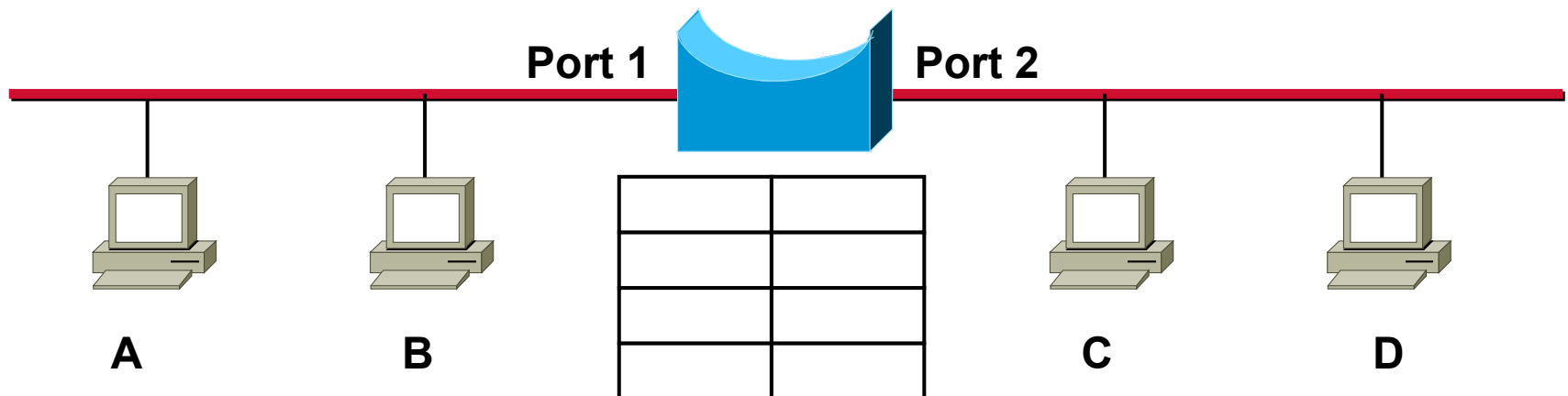


- **MAC addresses not routable**
 - ◆ NetBios over NetBEUI not routable (no L3)
- **Bridge supports different physical media on each port**
 - ◆ E.g. 10Mbit/s to 100Mbit/s
- **Router supports different layer-2 technologies**
 - ◆ E.g. Ethernet to Frame Relay

How does it work?

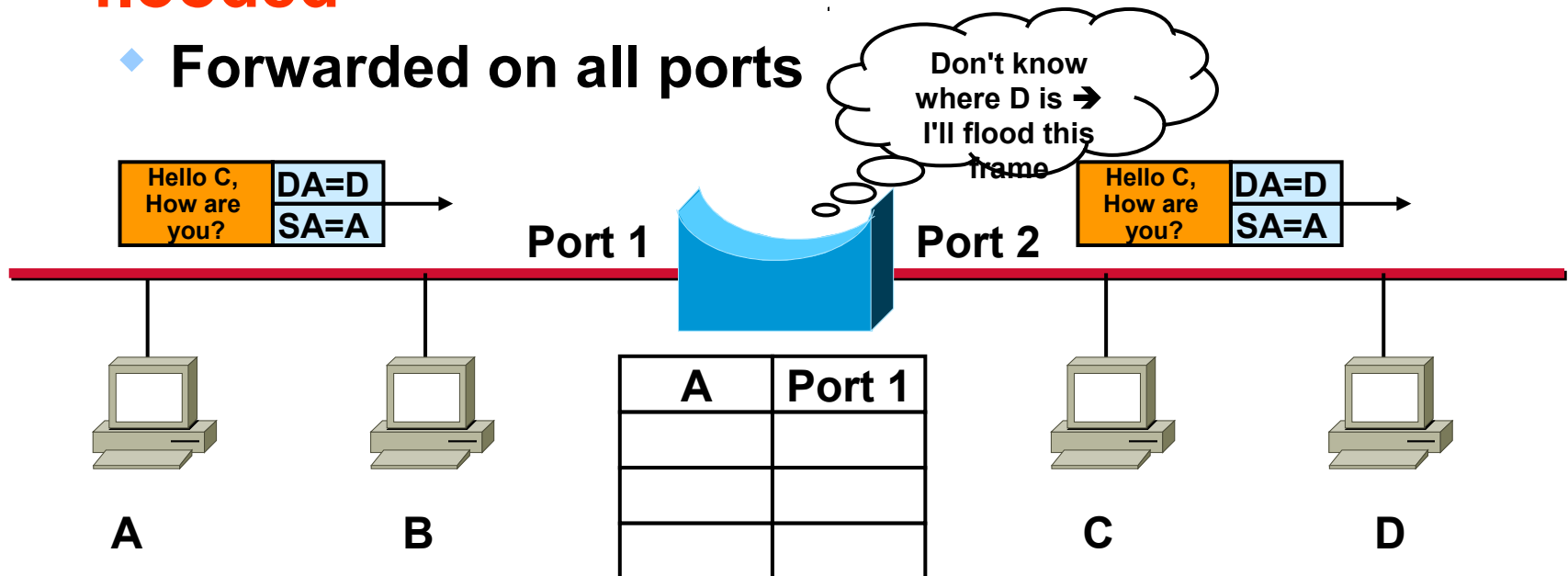


- Transparent bridging is like "plug & play"
- Upon startup a bridge knows nothing
- Bridge is in **learning mode**





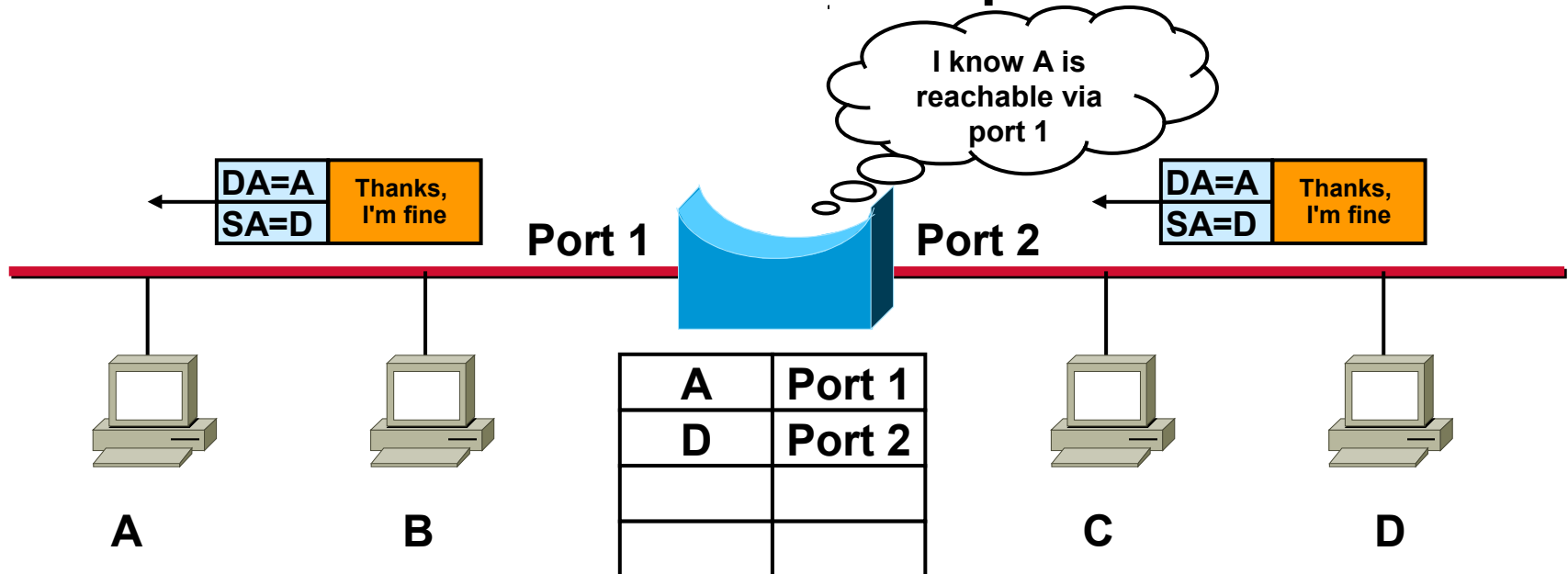
- Once stations send frames the bridge notices the **source** MAC address
 - ◆ Entered in bridging table
- Frames for unknown destinations are **flooded**
 - ◆ Forwarded on all ports



Learning → Table Filling



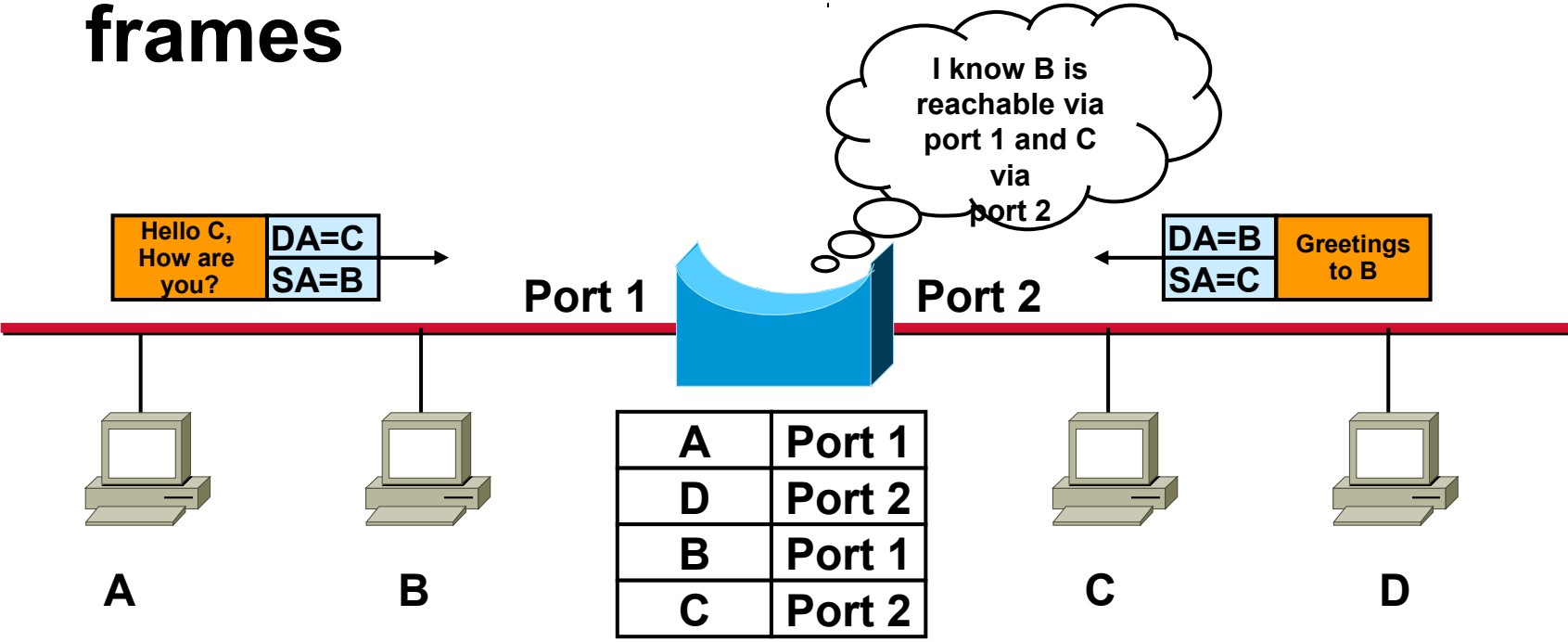
- If the destination address matches a bridging table entry, this frame can be actively
 - ◆ **forwarded** if reachable via other port
 - ◆ **filtered** if reachable on same port



Learning → Table Filling



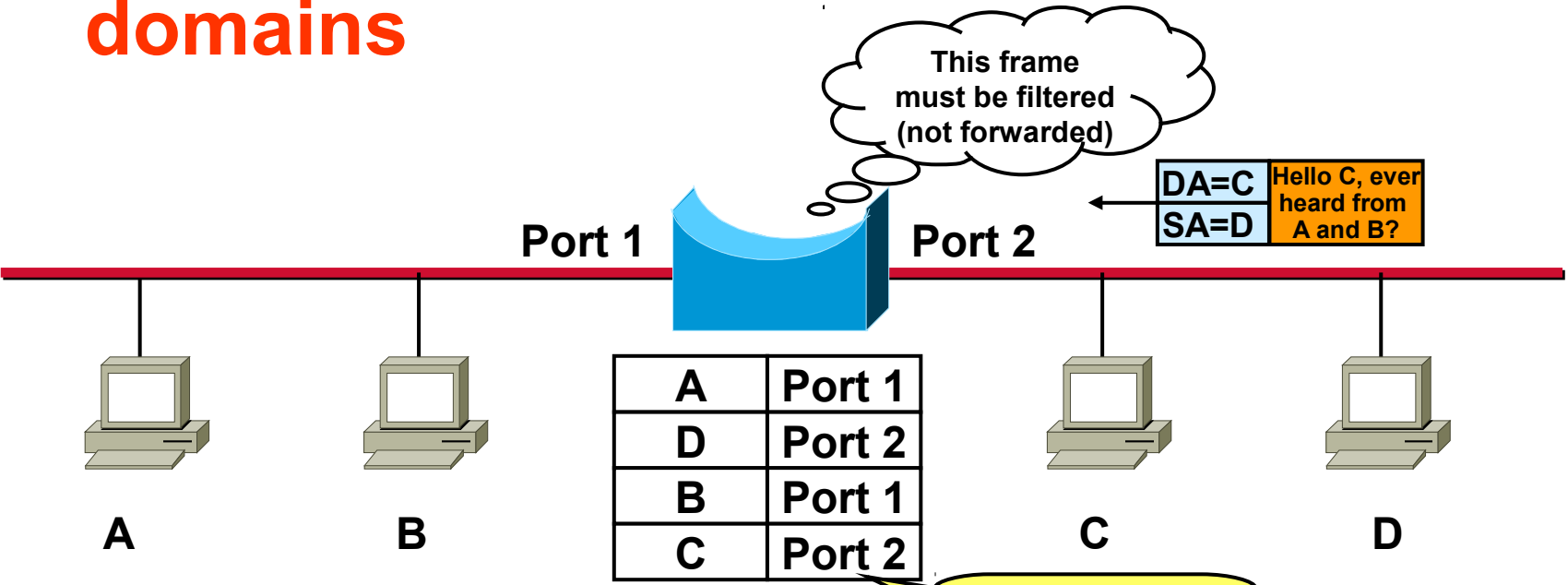
- After some time the location of every station is known – simply by listening!
- Now only **forwarding** and **filtering** of frames



Forwarding and Filtering



- Frames whose source and destination address are reachable over the same bridge port are filtered
- LAN separated into **two collision domains**



Most Important !

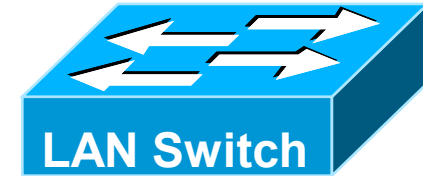


- Bridge separates LAN into **multiple collision domains** !
- A bridged network is still **one broadcast domain** !
 - ◆ Broadcast frames are always flooded
- A **router** separates the whole LAN into **multiple broadcast domains**

What is a Switch?



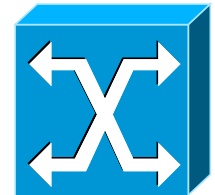
- A switch *is* basically a bridge, differences are only:

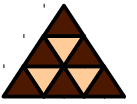


- ◆ **Faster** because implemented in **HW**
- ◆ **Multiple ports**
- ◆ **Improved functionality**

- **Don't confuse it with WAN Switching!**

- ◆ **Completely different !**
- ◆ **Connection oriented (stateful) VCs**





Bridge = Switch

Since we use only switches today, let's talk about them...

Modern Switching Features



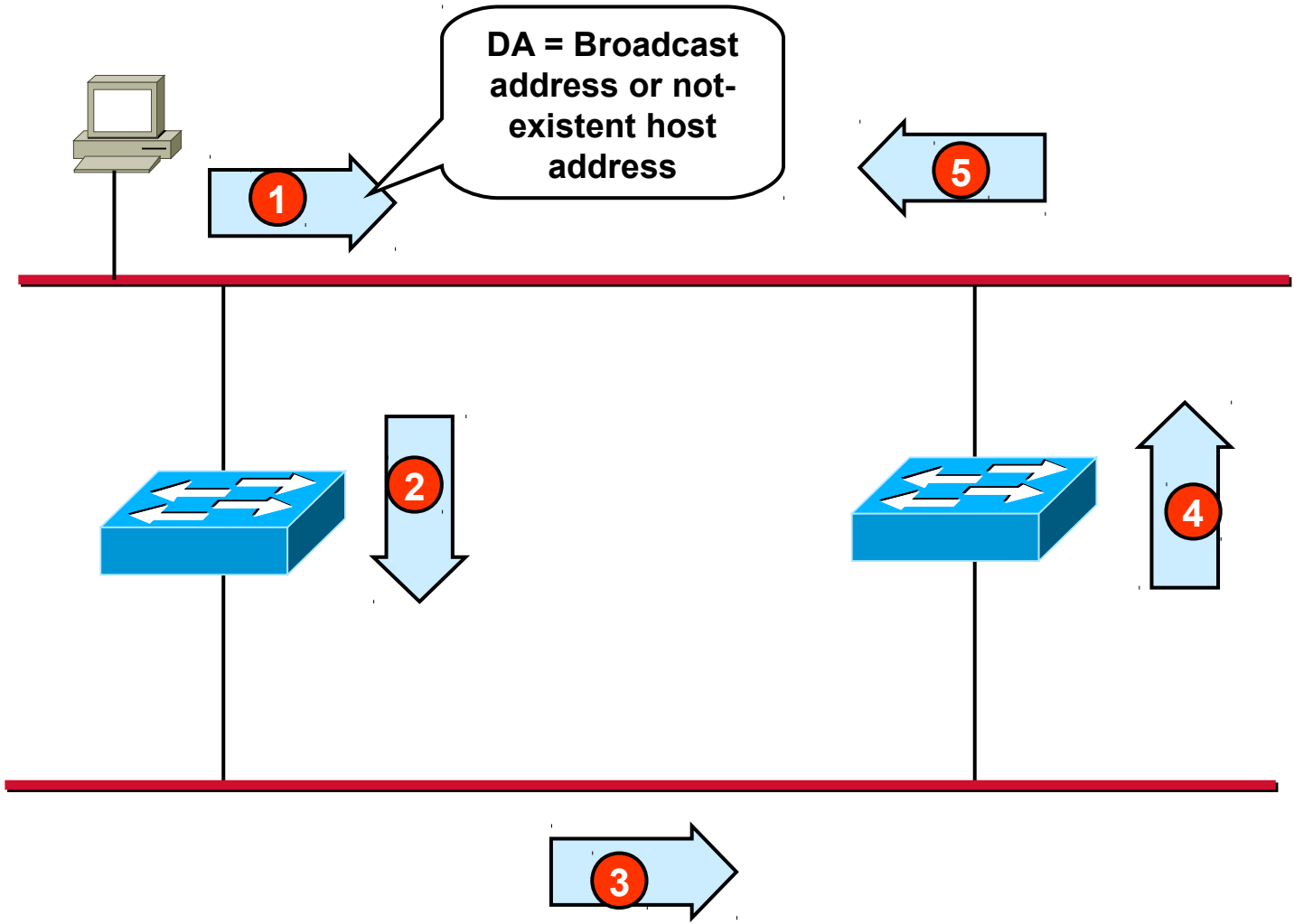
- **Different data rates supported simultaneously**
 - ◆ 10, 100, 1000, 10000 Mbit/s depending on switch
- **Full duplex operation**
- **QoS**
 - ◆ Queuing mechanisms
 - ◆ Flow control
- **Security features**
 - ◆ Restricted static mappings (DA associated with source port)
 - ◆ Port secure (Limited number of predefined users per port)
- **Different forwarding**
 - ◆ Store & Forward
 - ◆ Cut-through
 - ◆ Fragment-Free
- **VLAN support (Trunking)**
- **Spanning Tree**

Bridging Problems



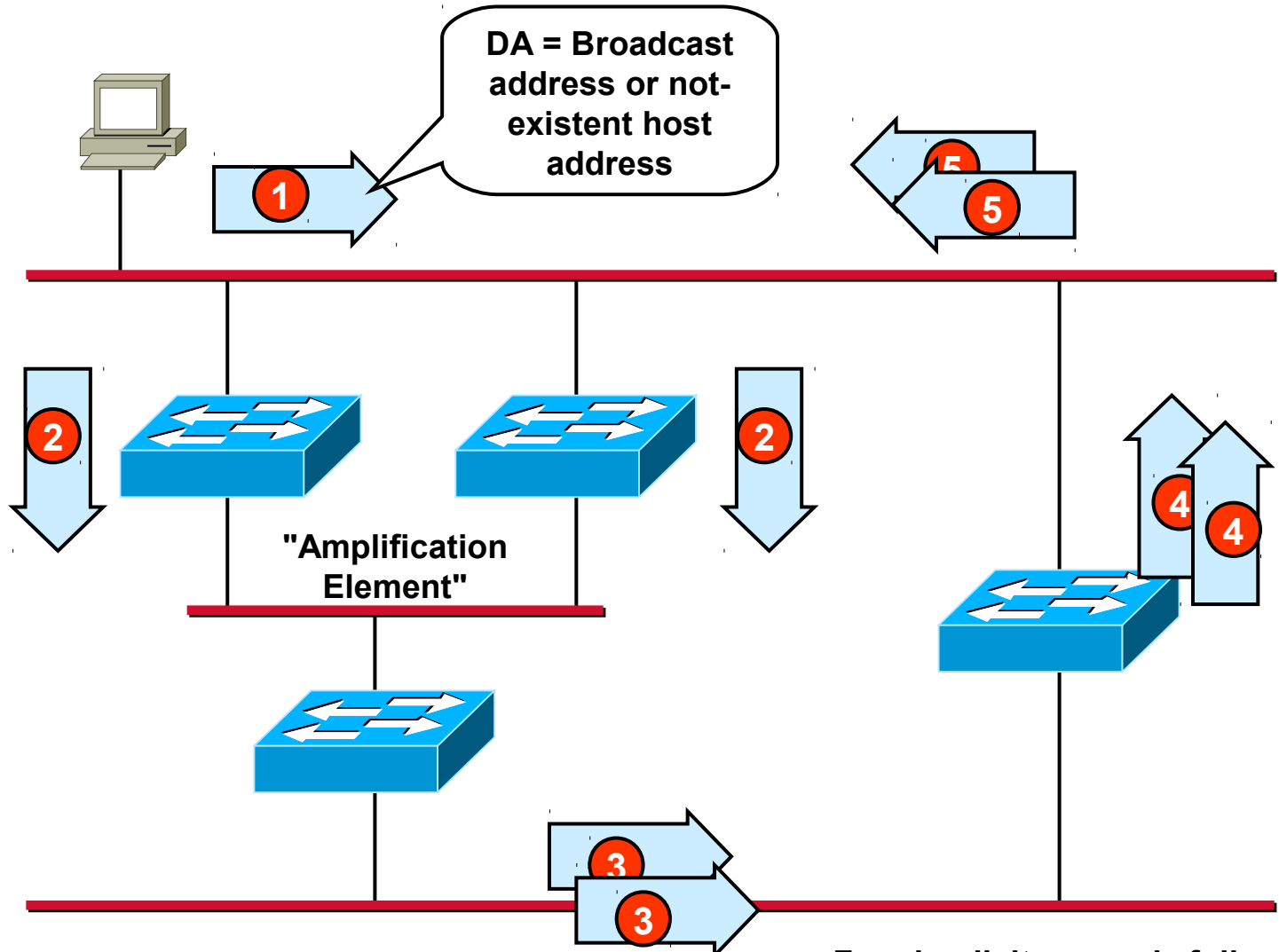
- **Redundant paths lead to**
 - ◆ **Broadcast storms**
 - ◆ **Endless cycling**
 - ◆ **Continuous table rewriting**
- **No load sharing possible**
- **No ability to select best path**
- **Frame may be stored for 4 seconds (!)**
 - ◆ **Although rare cases**
 - ◆ **But only little acceptance for realtime and isochronous traffic – might change!**

Endless Circling



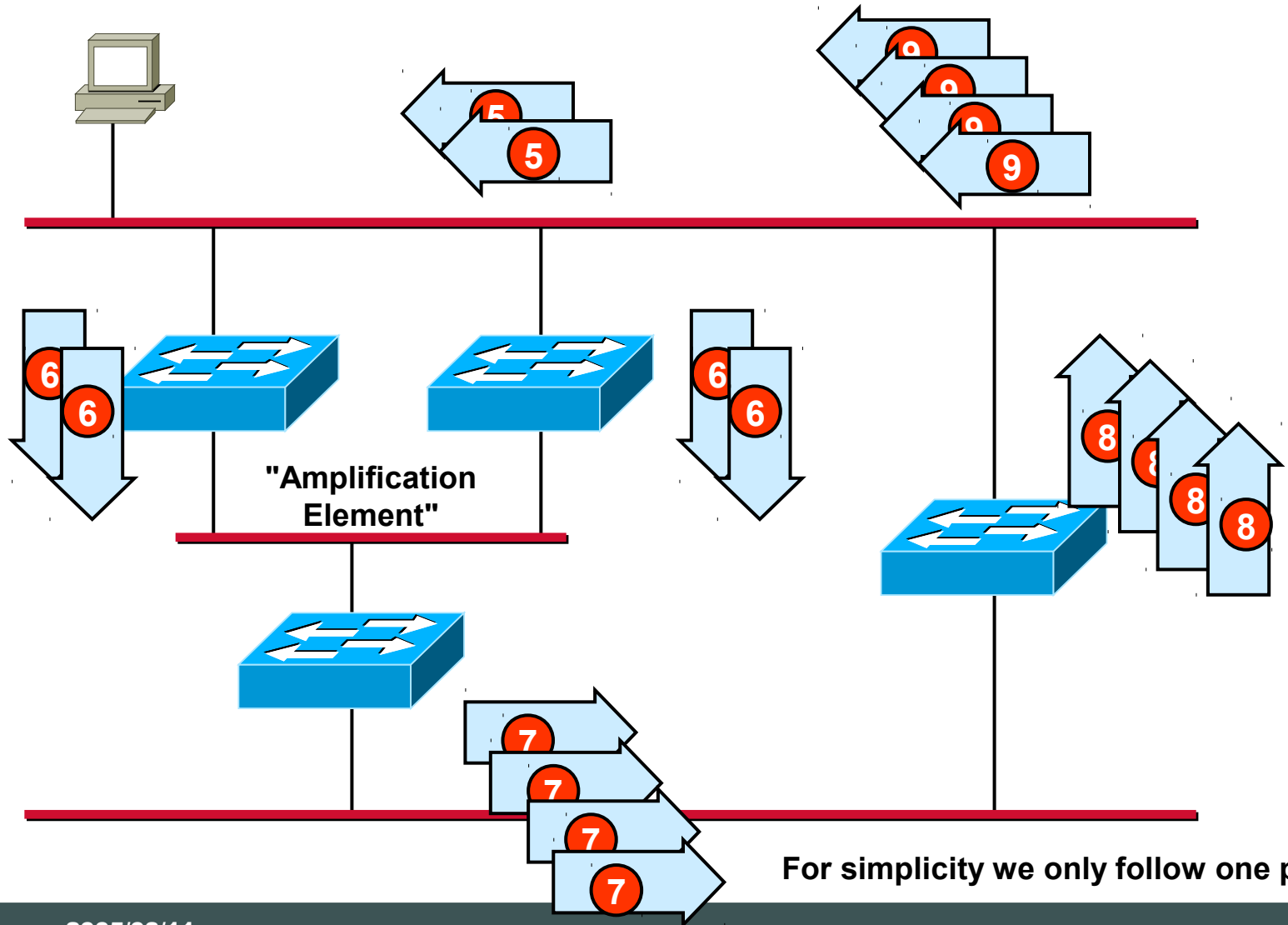
For simplicity we only follow one path

Broadcast Storm (1)

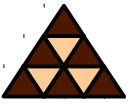


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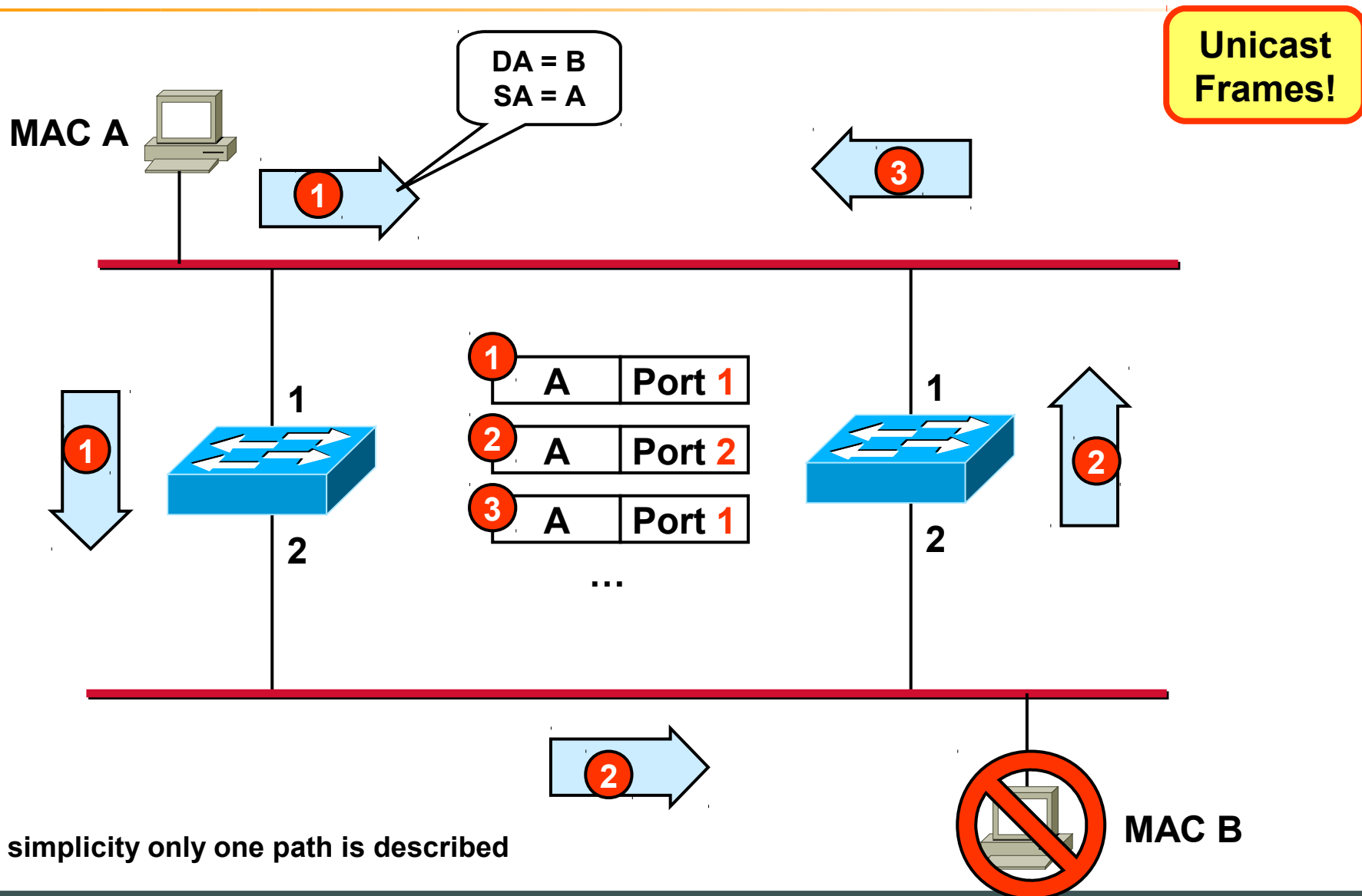
Broadcast Storm (2)



For simplicity we only follow one path



Mutual Table Rewriting



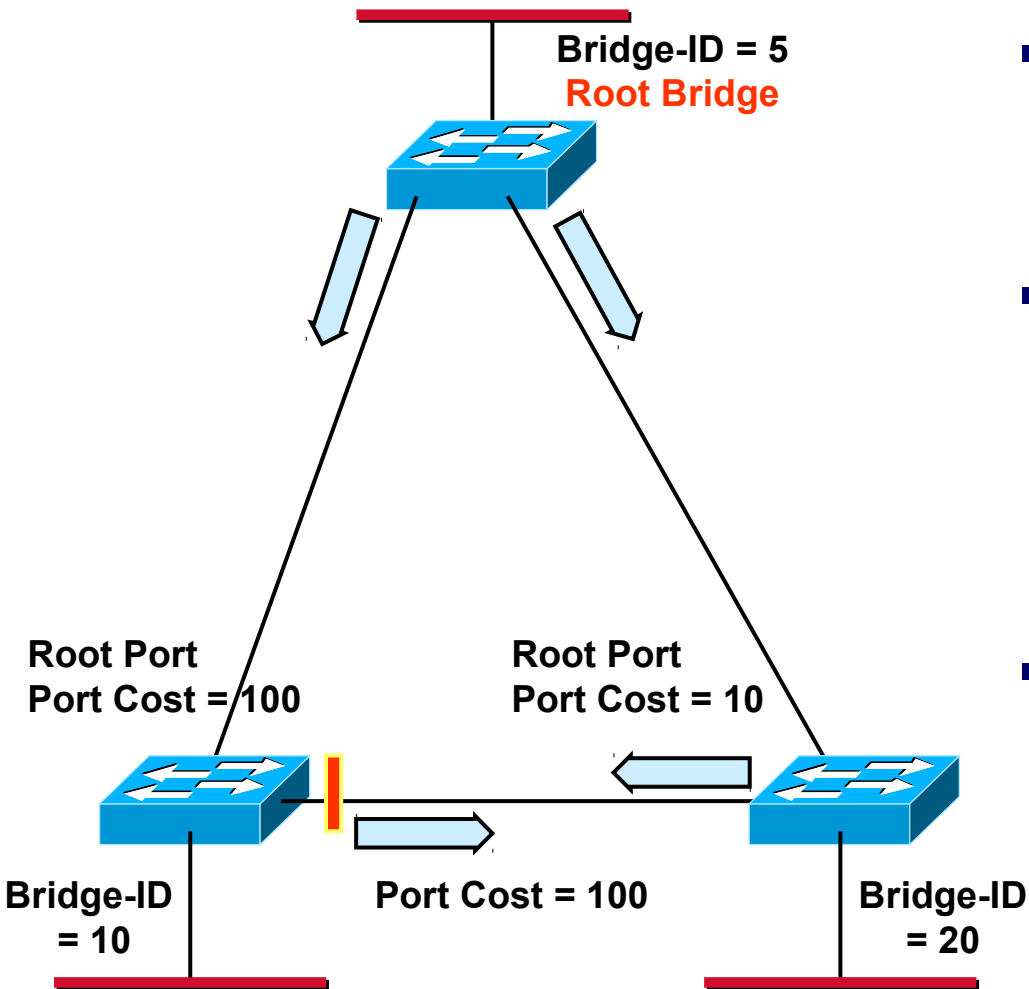


- Invented by *Radia Perlman* as general "mesh-to-tree" algorithm
- A must in bridged networks with redundant paths
- Only one purpose:
cut off redundant paths with highest costs

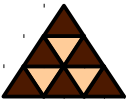


- **Special STP frames: "Bridge Protocol Data Units" (BPDU)**
- **A Bridge-ID for each bridge**
 - ◆ Priority value (16 bit, default 32768)
 - ◆ (Lowest) MAC address
- **A Port Cost for each port**
 - ◆ Default 1000/Mbits (can be changed)
 - ◆ E.g. 10 Mbit/s → C=100

STP Principle

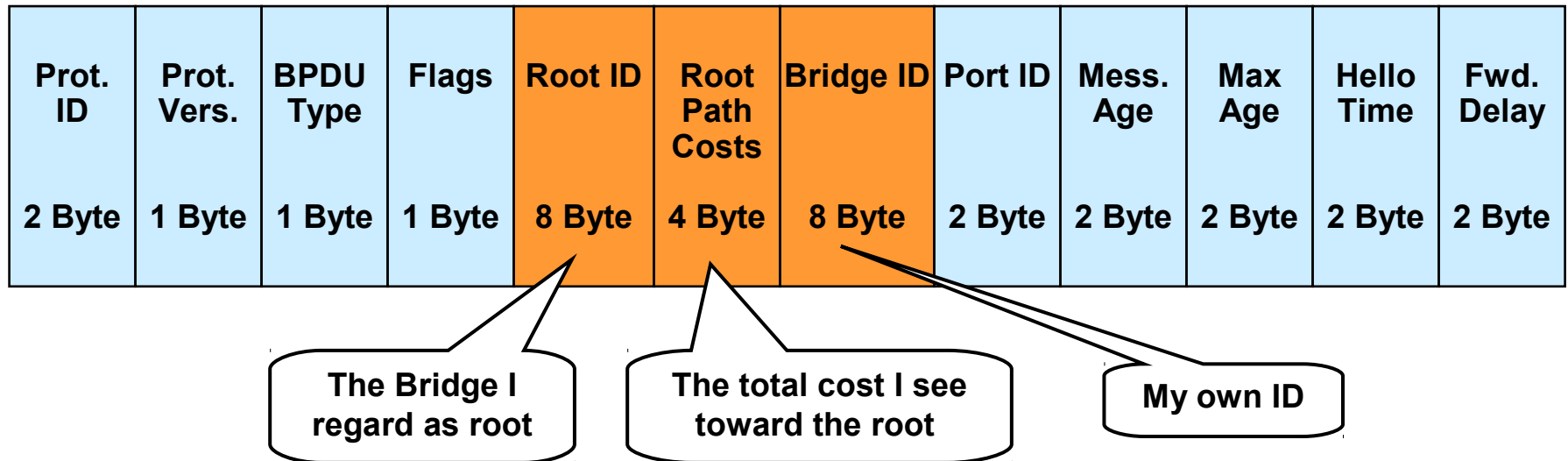


- First a **Root Bridge** is determined
 - ◆ Initially every bridge assumes itself as root
 - ◆ The bridge with lowest Bridge-ID wins
- Then the root bridge triggers BDPU sending (hello time intervals)
 - ◆ Received at "**Root Ports**" by other bridges
 - ◆ Every bridge adds its own port cost to the advertised cost and forwards the BDPU
- On each LAN segment one bridge becomes **Designated Bridge**
 - ◆ Having lowest total root path cost
 - ◆ Other bridges set redundant ports in **blocking state**



BPDU Format

- Each bridge sends periodically BPDUs carried in Ethernet multicast frames
 - ◆ Hello time default: 2 seconds
- Contains all information necessary for building Spanning Tree





- **Redundant links remain in active stand-by mode**
 - ◆ **If root port fails, other root port becomes active**
- **Low-price switches might not support STP**
 - ◆ **Don't use them in meshed configurations**
- **Only 7 bridges per path allowed according standard (!)**

Bridging versus Routing



Bridging



Depends on MAC addresses only



Invisible for end-systems;
transparent for higher layers



Must process every frame



Number of table-entries = number of all
devices in the whole network



Spanning Tree eliminates redundant lines;
no load balance



No flow control

Routing



Requires structured addresses (must be
configured)



End system must know its default-router



Processes only frames addressed to it



Number of table-entries = number of
subnets only



Redundant lines
and load balance possible








Flow control is possible
(router is seen by end systems)






Bridging versus Routing



Bridging

-  No LAN/WAN coupling because of high traffic (broadcast domain!)
-  Paths selected by STP may not match communication behaviour/needs of end systems
-  Faster, because implemented in HW; no address resolution
-  Location change of an end-system does not require updating any addresses
-  Spanning tree necessary against endless circling of frames and broadcast storms

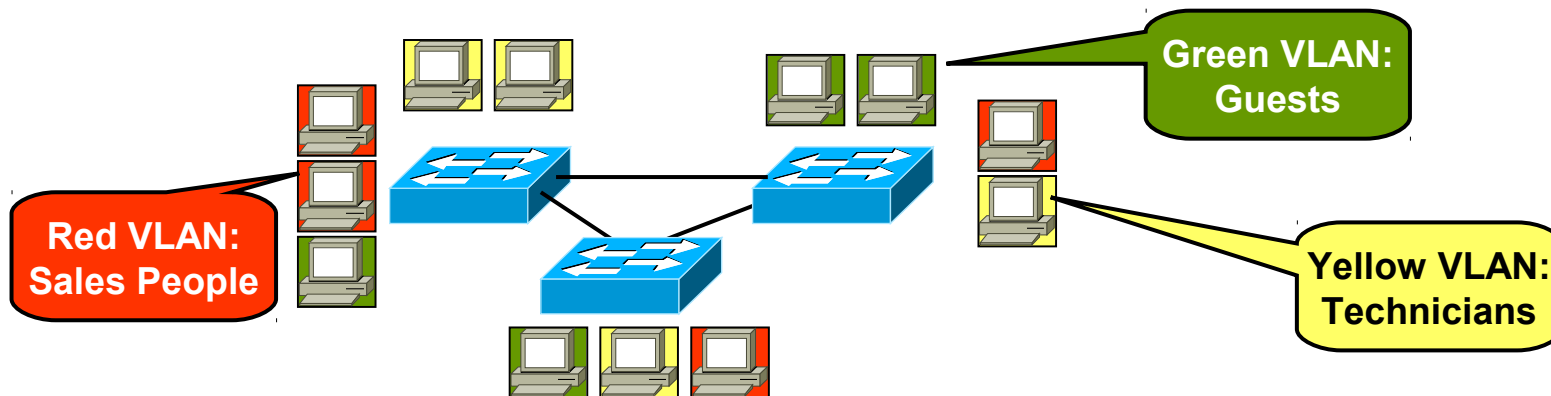
Routing

-  Does not stress WAN with subnet's broadcast or multicasts; commonly used as "gateway"
-  Router knows best way for each frame
-  Slower, because usually implemented in SW; address resolution (ARP) necessary
-  Location change of an end-system requires adjustment of layer 3 address
-  Routing-protocols necessary to determine network topology



Virtual LANs

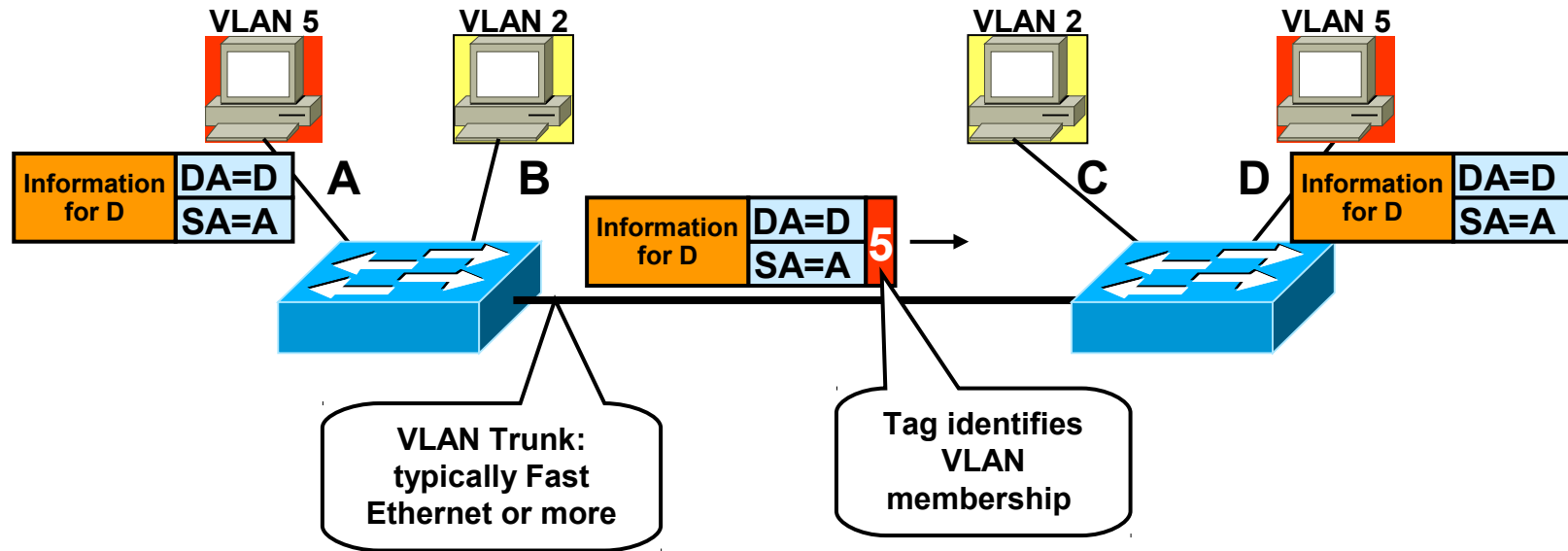
- **Separate LAN into multiple broadcast domains**
 - ◆ No global broadcasts anymore
 - ◆ For security reasons
- **Assign users to "VLANs"**





- **Different solutions**
 - ◆ **Port** based assignment
 - ◆ **Source address** assignment
 - ◆ Protocol based
 - ◆ Complex rule based
- **Bridges are interconnected via **VLAN trunks****
 - ◆ **IEEE 802.1q** (New: 802.1w, 802.1s)
 - ◆ **ISL (Cisco)**

VLAN Trunking Example

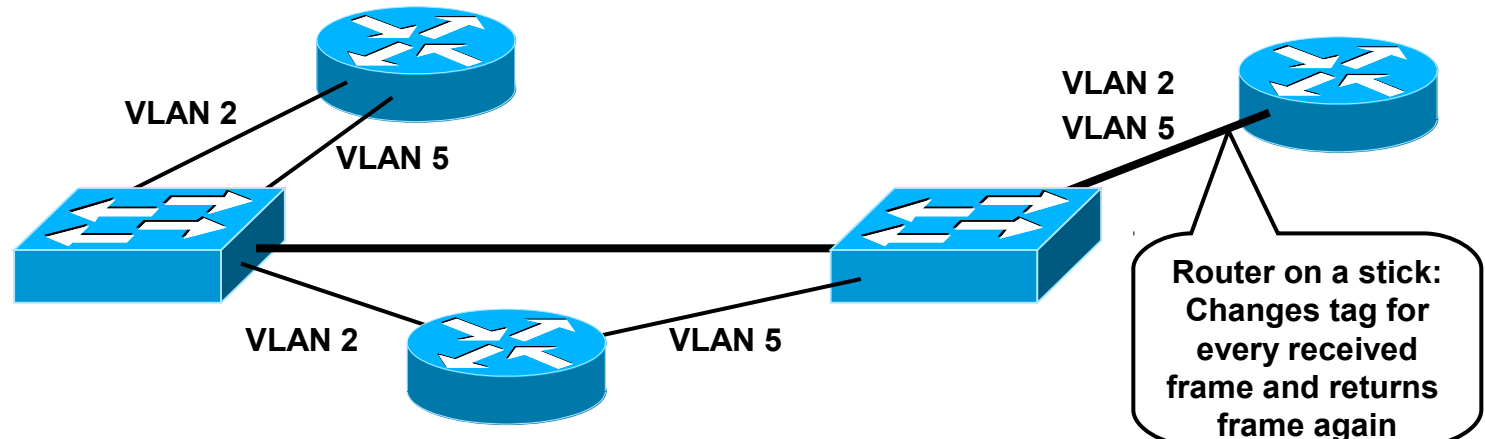


- Inter-VLAN communication not possible
- Packets across the VLAN trunk are **tagged**
 - ◆ Either using 802.1q or ISL tag
 - ◆ So next bridge is able to constrain frame to same VLAN as the source

Inter-VLAN Traffic



- Router can forward inter-VLAN traffic
 - ◆ Terminates Ethernet links
 - ◆ Requirement: **Each VLAN in other IP subnet !**
- Two possibilities
 - ◆ Router is member of every VLAN with one link each
 - ◆ Router attached on VLAN trunk port ("Router on a stick")





- Ethernet Bridging is "**Transparent Bridging**"
 - ◆ Hosts do not "see" bridges
 - ◆ Plug & Play
- **1 Collision domain → 1 Broadcast domain**
- Switches increase network **performance !**
- Redundant paths are dangerous
 - ◆ Broadcast storm is most feared
 - ◆ Solution: **Spanning Tree Protocol**
- **VLANs create separated broadcast domains**
 - ◆ Port based or address based VLANing
 - ◆ Routers allow inter-VLAN traffic



- **Can I bridge from Ethernet to Token Ring?**
- **How is flow control implemented?**
- **Which bridge should be root bridge?**
- **What are main differences between 802.1q and ISL?**
- **What are Layer-3, Layer-4, and Layer-7 switches ?**